

Lecture outline



- Basic inventory problems
- The economic order quantity
- An inventory game
- Multiperiod lot sizing
 - » Math programming formulation
 - » Heuristics
 - » Wagner-Whitin algorithm

Basic inventory problems

■ Examples:

» Products:

- Customers consume products over time.
- Store replenishes periodically.

» People with specialized training:

- People randomly leave the company over time.
- Company periodically hires new graduates.

» Water (management of dams).

- Rainfall randomly replenishes reservoirs.
- Release water from dam to maintain level.

» Oil being stored in storage tanks:

- Oil is steadily consumed.
- Periodically is replenished from tankers.

Basic inventory problems

■ Examples:

- » Housing stock
 - Houses are continually being purchased.
 - Developers produce new developments or apartment buildings.
- » Financial resources (startup company)
 - Cash is used to build up the company.
 - Periodically fresh capital is raised from venture capitalists.
- » Features in a software program:
 - Accumulating features in a software program in response to user requests and the ideas of developers, or due to bug fixes.
 - Periodically ship a new version of the program.
- » Purchasing stock:
 - Funds become available for investment.
 - Periodically purchase new shares of stock.

Basic inventory problems

Statement Date	Account Number	Account Summary for the Period	Replenishment Amount	Replenishment Method
1/4/2002		11/4/2001 to 1/3/2002	25.00	AMEX

Tag Deposit	Beginning Balance	Tolls & Fees	Payments & Credits	Ending Balance	Replenishment Threshold
0.00	22.07	27.25	25.00	19.82	10.00

Date/Time	Tag	Transaction	Entry		Exit		Class	Amount	Balance
			Plaza	Lane	Plaza	Lane			
11/03 14:20	02200917779	New Jersey Turnpike Toll	9	05E	13A	13X	1	-1.55	20.52
11/06 22:21	02200917779	New Jersey Turnpike Toll	13A	05E	9	13X	1	-1.45	19.07
11/11 15:09	02200917779	New York State Thruway Toll	15	05E	24	08S	1	-3.65	15.42
11/13 14:27	02200917779	New York State Thruway Toll	24	04E	15	08W	1	-3.65	11.77
11/13 17:35	02200917779	New Jersey Turnpike Toll	10	11E	7A	06X	1	-0.75	11.02
12/15 16:59	02200917779	New Jersey Turnpike Toll	9	05E	16E	01X	1	-2.30	8.72
12/15 17:42	02200917779	PANYNJ Toll			LT	01	1	-5.00	3.72
12/16 10:53	02200917779	New Jersey Turnpike Toll	14C	08E	9	12X	1	-2.40	1.32
12/17 16:50	02200917779	New Jersey Turnpike Toll	9	05E	13A	13X	1	-1.55	-0.23
12/18 02:35		Replenishment						25.00	24.77
12/19 20:53	02200917779	New Jersey Turnpike Toll	13A	06E	9	13X	1	-1.45	23.32
12/26 13:44	02200917779	New York State Thruway Toll	15	02E	23	06E	1	-3.50	19.82

Basic inventory problems

■ Mutual fund cash balance

Stock market



Cash



Investor



How much cash do we keep on hand to strike a balance between the deposits and withdrawals of investors, and the behavior of the market?

Basic inventory problems

- There are a number of ways to refer to storing resources for the future:
 - » Physical resources
 - Inventory
 - Stockpile
 - Stock

 - » Financial resources
 - Savings
 - Nest egg
 - Reserve

Basic inventory problems

■ Reasons for holding inventories:

- » Economies of scale
 - Batches of goods
 - Discounts (purchasing)
 - Transportation economies (e.g. shipping in bulk)
- » Uncertainties
 - Demand
 - Order lead times
 - Supply/price of raw materials (OPEC)
 - Supply/price of components (strikes)
 - Quality control
- » Speculation
 - commodities prices
 - currency fluctuations

Basic inventory problems

■ Reasons for holding inventories

» Transportation

- In-transit or pipeline inventories

» Smoothing production

- Respond to seasonal patterns in demand
- Seasonal production of some items
 - Certain foods
 - Syrup
 - Snow
 - Students

» Control costs

- Lower inventories requires more sophisticated control systems

Basic inventory problems

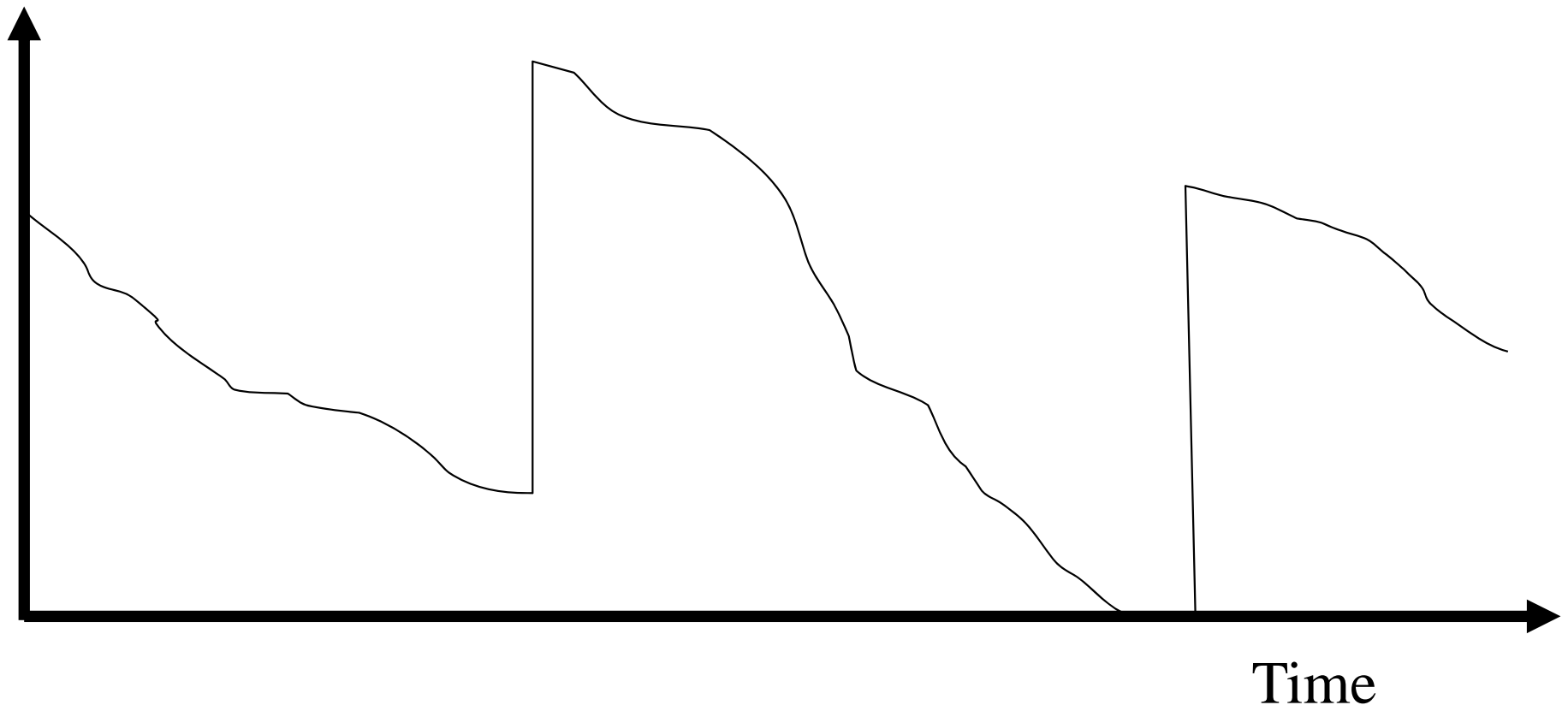
■ The lot sizing problem

- » Often, there are economies of scale when ordering new resources:
 - Raising operating capital
 - There is a fixed cost to going to the capital markets
 - Just as much work to raise \$1m as \$5m
 - Shipping the latest version of a software program
 - New features are added over time
 - There is a fixed cost of shipping a new version of the code
 - How many new features do you add before you ship the code?
 - Ordering new product for a store shelf
 - Fixed cost for placing and shipping an order

Basic inventory problems

- An aging and replenishment process (negative drift):

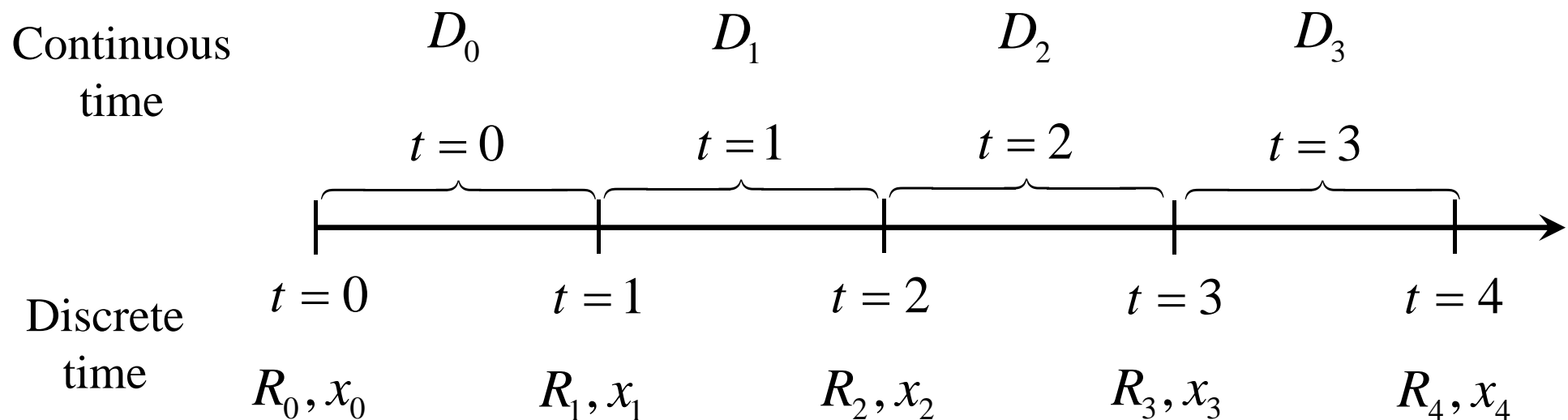
- » State is inventory – drift is due to customer demand.



Basic inventory problems

■ Indexing time:

- » Deterministic indexing – Index based on when something happens.
- » Stochastic indexing – Index based on when something becomes known.
- » Deterministic indexing:

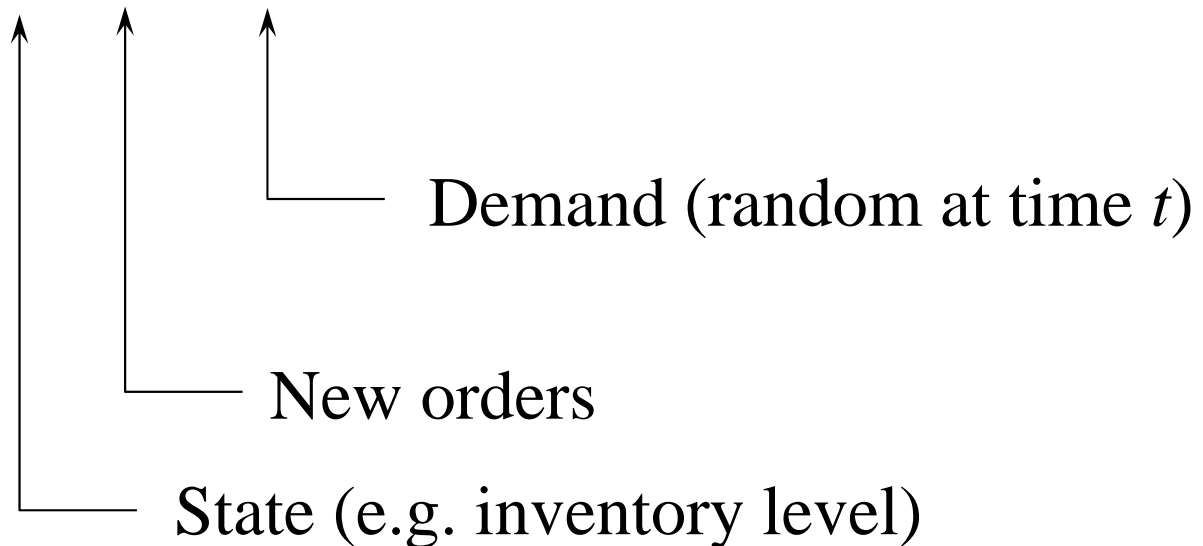


Basic inventory problems

■ Basic inventory equation

» Deterministic indexing

$$R_{t+1} = \max \{0, R_t + x_t - D_t\} = [R_t + x_t - D_t]^+$$

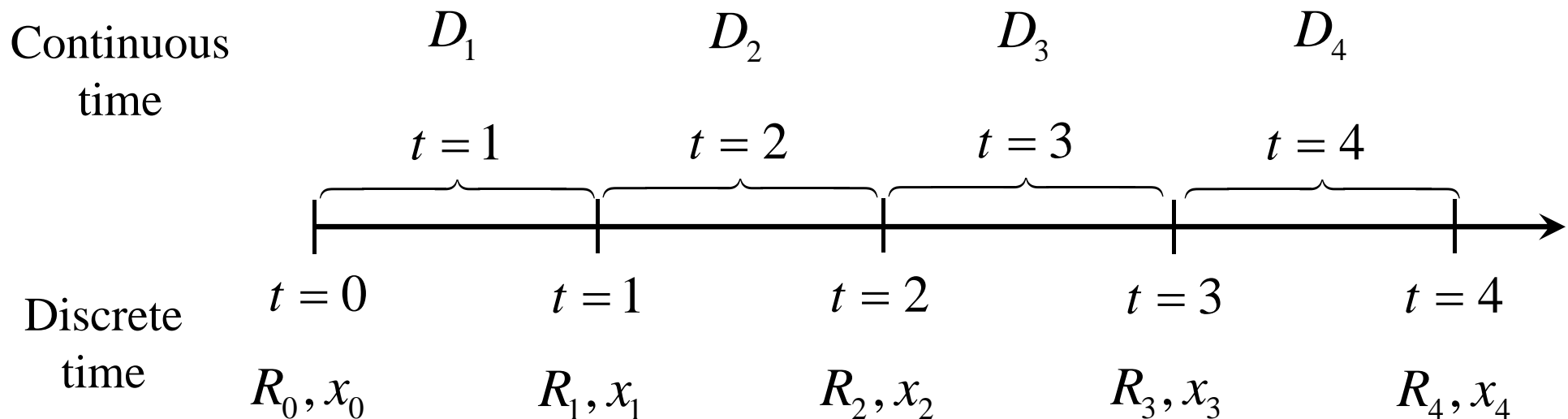


» In deterministic indexing, everything is modeled at the beginning of a time period.

Basic inventory problems

■ Stochastic indexing

- » Information arrives continuously over time
- » *A variable indexed by t contains exogenous information up through time t .*



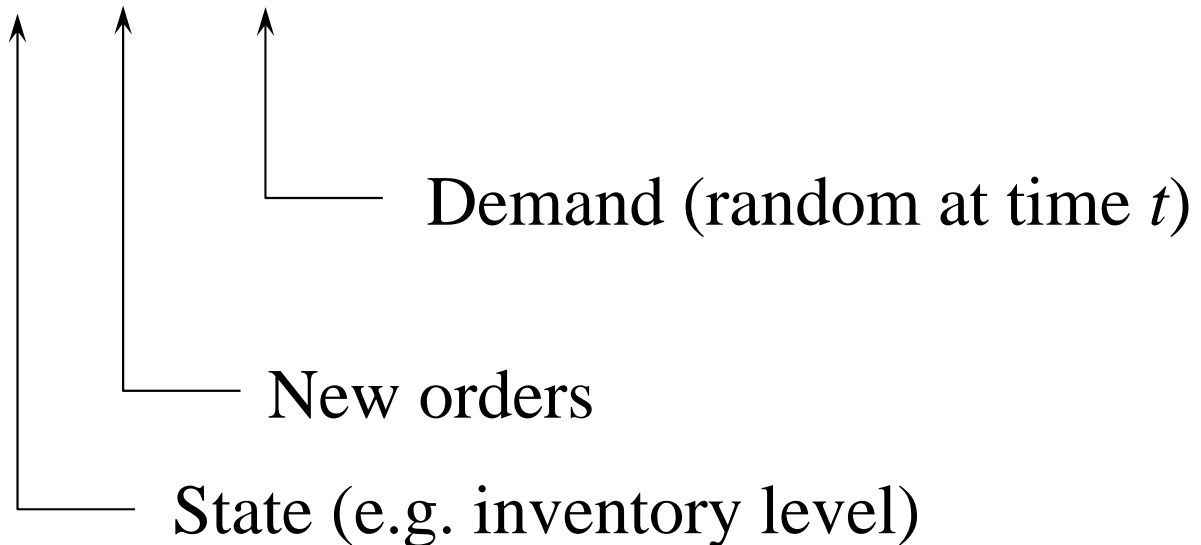
- » In stochastic indexing, everything is indexed at the end of a period.

Basic inventory problems

■ Basic inventory equation

» Stochastic indexing

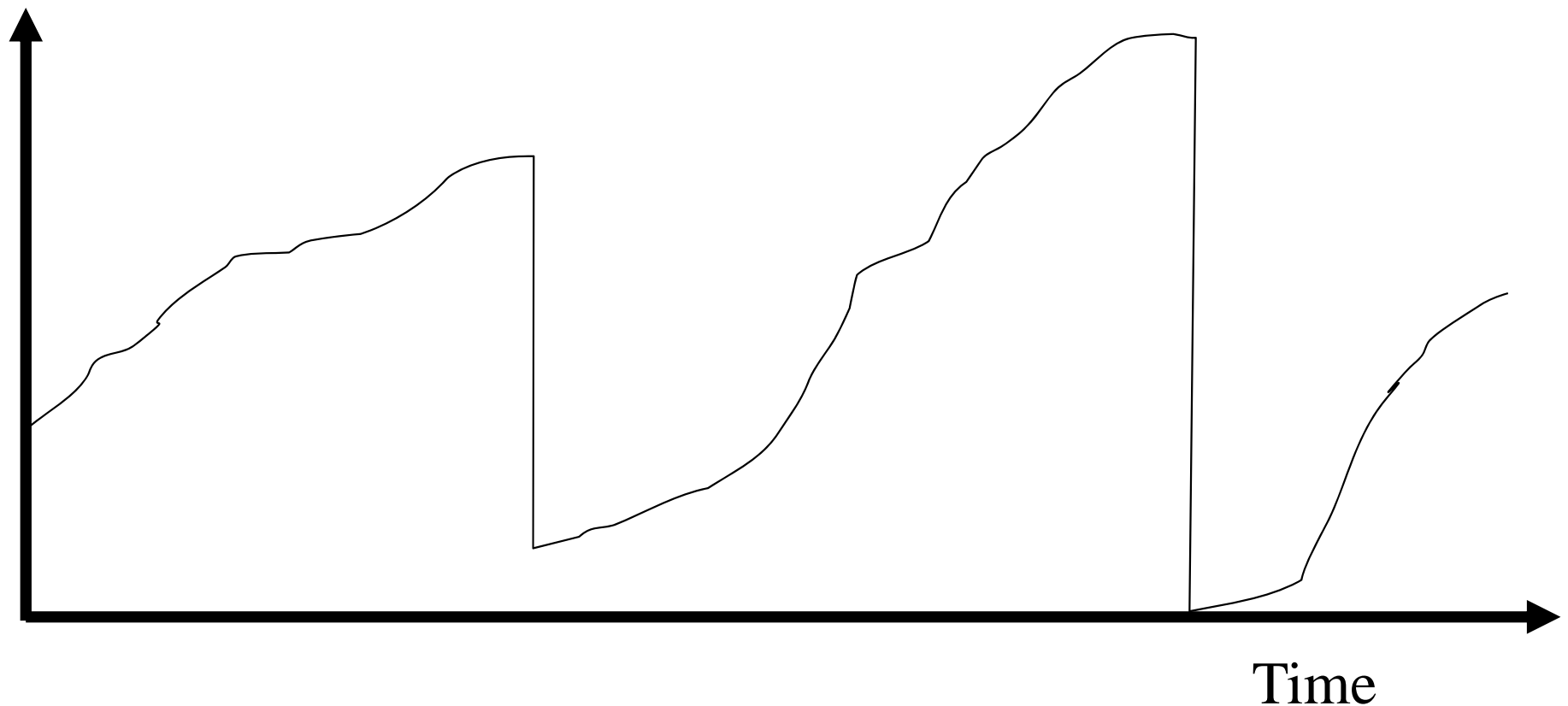
$$R_{t+1} = \max \{0, R_t + x_t - D_{t+1}\} = [R_t + x_t - D_{t+1}]^+$$



Basic inventory problems

■ Inventory with positive drift

- » Cash is deposited in an account. Periodically the cash is invested in batch amounts to reduce transaction costs.



Basic inventory problems

- Basic inventory with positive drift
 - » Blood donations, water reservoirs, ...

$$R_{t+1} = [R_t - x_t]^+ + D_{t+1}$$

State (e.g. accumulation)

Endogenously controlled shift

Exogenous drift

Outline



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 - » Wagner-Whitin algorithm

The economic order quantity

■ Order costs

- » Fixed cost of placing an order
 - Paperwork, forms, telephone
 - Sending a truck out
 - Setting up a machine
 - Price structure imposed by supplier
- » Variable cost of ordering a certain amount of product
 - Variable cost may be fixed (a linear function).
 - Or there may be economies of placing larger orders.

The economic order quantity

■ Holding cost

- » Storage (heat, electricity, supervision, etc.)
- » Taxes and insurance
- » Breakage, spoilage, deterioration and obsolescence
 - Careful with these - these “costs” convert to reduction in quantity.
- » Opportunity cost (interest)
 - Hurdle rate for a company is generally much higher than bank interest rates.
 - Let I = “interest rate”
 - c^p = purchase cost of item
 - $c^h = Ic^p$ = holding cost (be careful with units; if I is interest rate per year, c^h is holding cost per year).

The economic order quantity

■ Stockout cost

- » Cost of lost customers
- » Cost of pushing orders to future time periods

■ Notes:

- » Stockout costs are not relevant in our simplest inventory system, because they cannot happen.
- » Stockouts arise when:
 - Demand is random
 - Demand varies over time with production capacities
 - Order costs may vary as a function of time, possibly exceeding the “benefit” of covering demand.

The economic order quantity

■ Some assumptions:

- » Demand is deterministic with rate λ per unit time.
- » Rate λ is constant – stationary process.
- » Costs are stationary.
- » Orders arrive immediately.
- » All orders must be filled.

The economic order quantity

The basic inventory equation:

$$R_{t+1} = [R_t + x_t - D_{t+1}]^+$$

where:

R_t = Inventory at start of time t

x_t = Amount ordered at time t

D_t = Demand during period t

Notation:

$$[x]^+ = \max \{x, 0\}$$

The economic order quantity

■ The cost function:

$$\begin{aligned}c(x_t, R_t) &= \text{Total costs during period } t \text{ given order quantity } x_t \text{ and initial} \\ &\quad \text{inventory } R_t \\ &= c^o(x_t) + c^h(R_t, x_t)\end{aligned}$$

where:

$$c^o(x_t) = \text{Order costs}$$

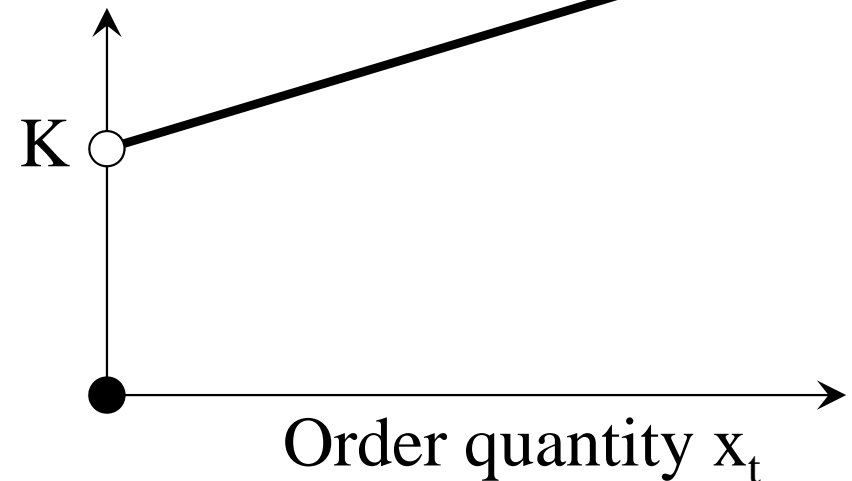
$$= \begin{cases} K + c^p x_t & x_t > 0 \\ 0 & x_t = 0 \end{cases}$$

$$c^p = \text{Unit purchasing costs}$$

$$c^h(R_t, x_t) = \text{Holding costs}$$

$$= c^h \int_0^{\Delta t} [R_t + x_t - \lambda z]^+ dz$$

$$c^h = \text{Unit holding costs per time}$$



The economic order quantity

■ The cost function:

We can transform the nonlinear cost function into a linear one:

$$c^o(x_t) = Ky_t + c^p x_t$$

where:

$$x_t \geq 0$$

$$x_t \leq My_t \quad M = \text{"big M"}$$

$$y_t \in (0,1)$$

If $y_t = 0$ then we force $x_t = 0$. Now we have transformed a nonlinear cost function into a linear one, but we have added an integer variable.

The economic order quantity

■ Infinite horizon problem:

We would like to solve:

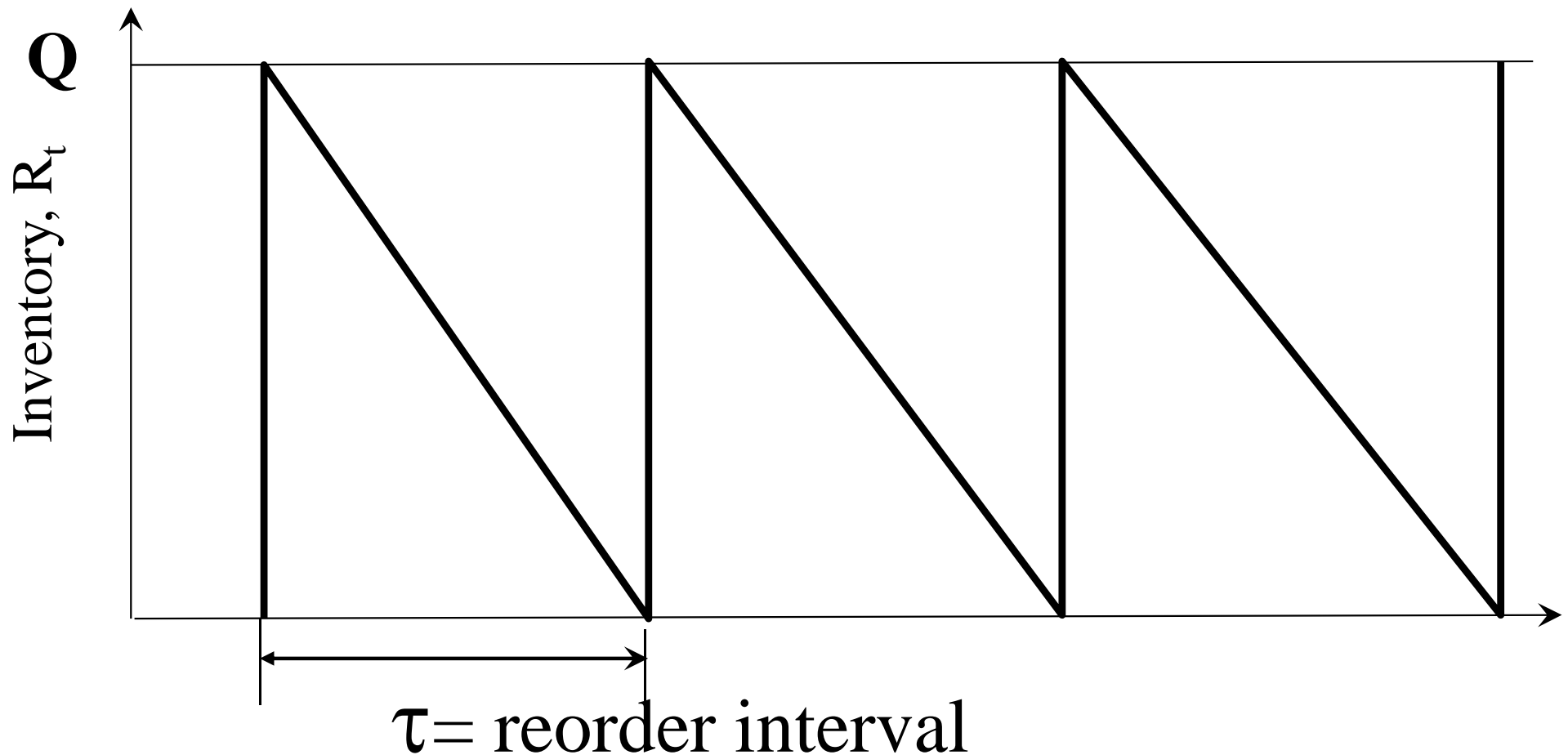
$$\min_{x_t, y_t} \sum_{t=0}^{\infty} c(R_t, x_t, y_t)$$

This is a really big number! A more formal way to write it is as an average cost:

$$\lim_{T \rightarrow \infty} \frac{1}{T} \left\{ \min_{x_t, y_t} \sum_{t=0}^T c(R_t, x_t, y_t) \right\}$$

The economic order quantity

Intuition suggests that we let inventories drop to zero, and then “order up to” an amount Q :



The economic order quantity

Reformulate the problem in terms of cost per order interval :

Q = order quantity

λ = Demand rate per unit time (assumed constant and deterministic)

$c^o(Q)$ = Order cost per order interval

= K (since we always make one order per interval).

τ = Length of order interval

$$= \frac{Q}{\lambda}$$

$c^p(Q)$ = Purchase costs per order interval

$$= c^p Q$$

The economic order quantity

$$\begin{aligned}c^h(Q) &= \text{Holding cost per order interval} \\ &= (\text{holding cost}) * (\text{average inventory}) * (\text{length of interval}) \\ &= c^h \left(\frac{Q}{2} \right) \tau \\ &= c^h \left(\frac{Q}{2} \right) \left(\frac{Q}{\lambda} \right) = c^h \left(\frac{Q^2}{2\lambda} \right)\end{aligned}$$

$$\begin{aligned}c(Q) &= \text{Total cost per order interval} \\ &= c^o(Q) + c^p(Q) + c^h(Q) \\ &= K + c^p Q + c^h \left(\frac{Q^2}{2\lambda} \right)\end{aligned}$$

The economic order quantity

How do we find the optimum value of Q ?

We can try differentiating $C(Q)$ with respect to Q :

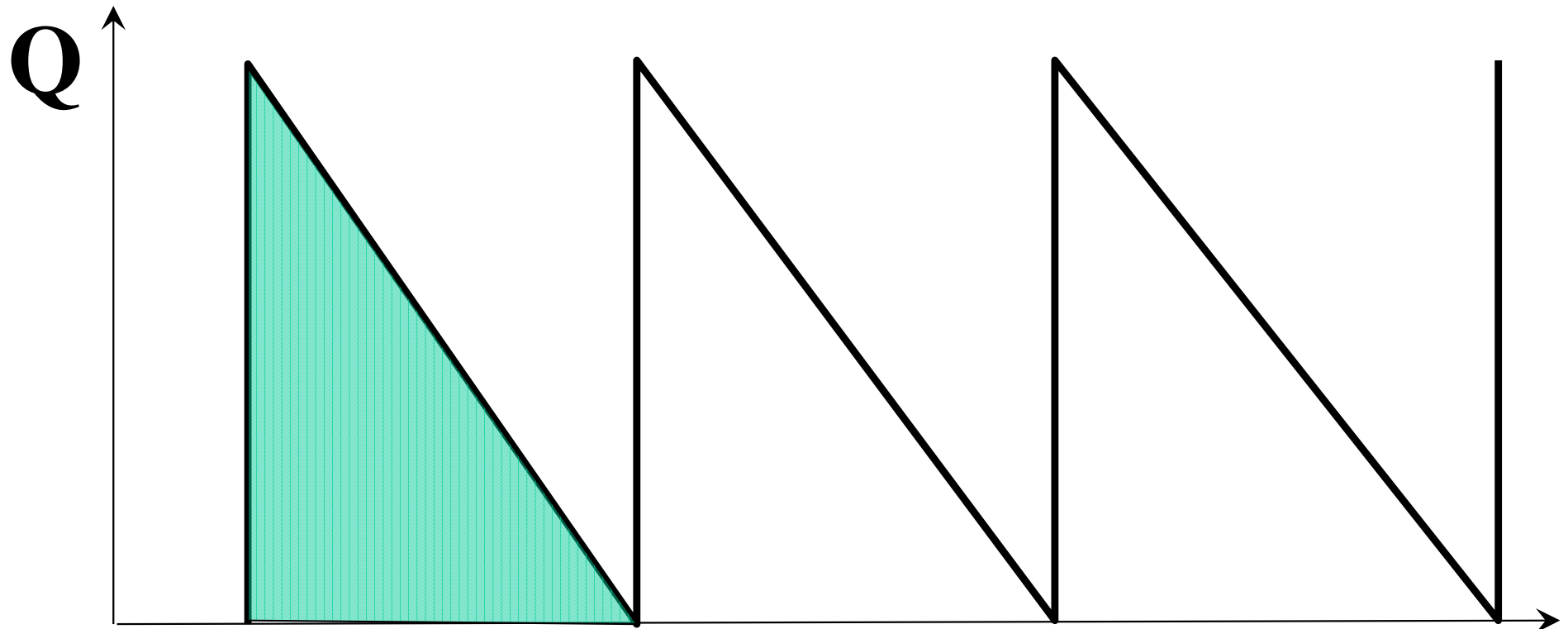
$$\frac{dC(Q)}{dQ} = 0 + c^p + 2c^h \frac{Q}{2\lambda} = 0$$

Solving for Q gives us:

$$Q^* = -\frac{c^p \lambda}{c^h} \quad !!!!!$$

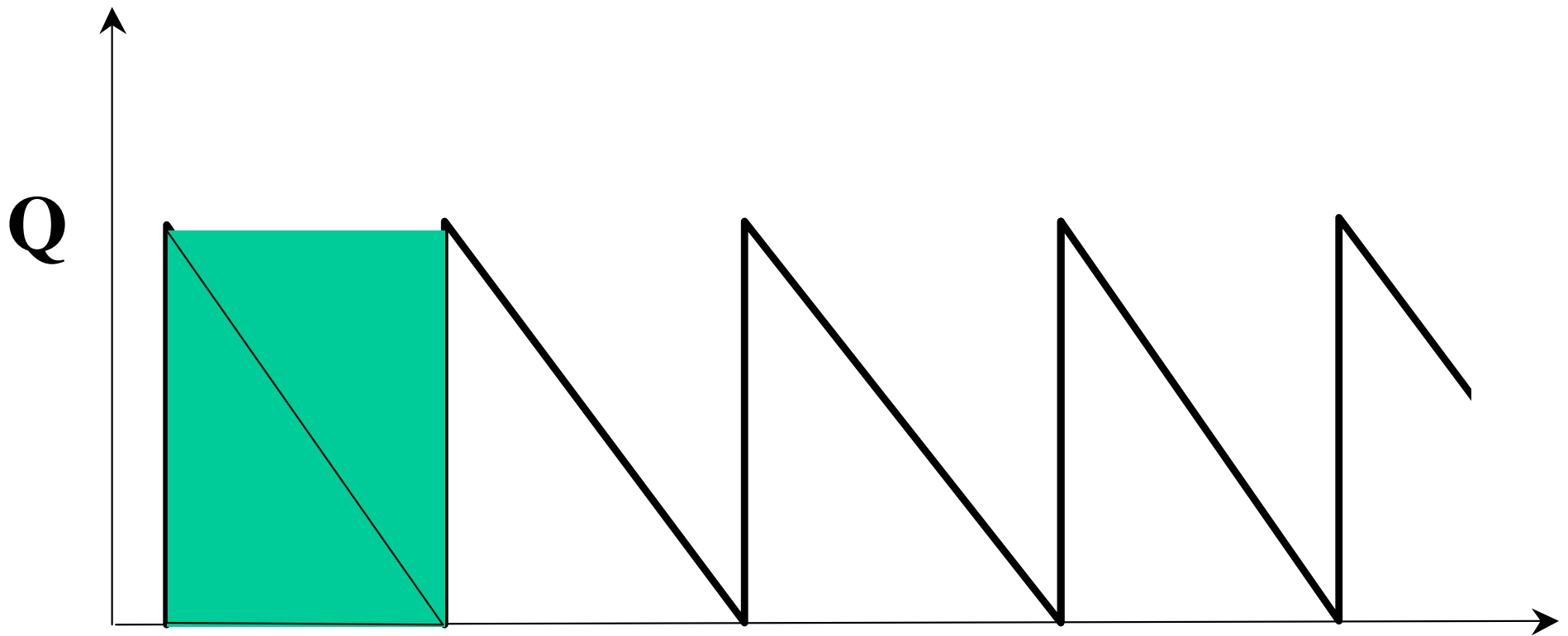
What went wrong?

The economic order quantity

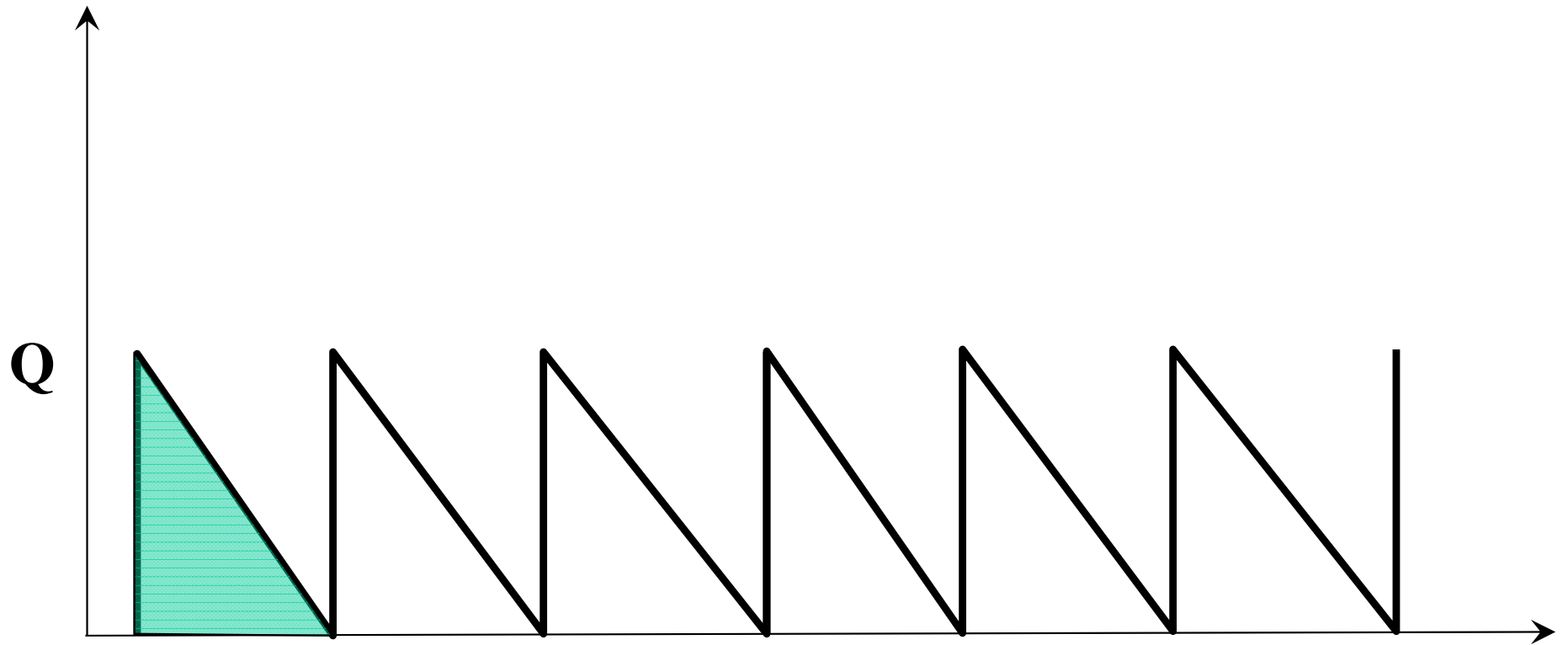


**Cost per cycle = K plus quantity
proportional to green area**

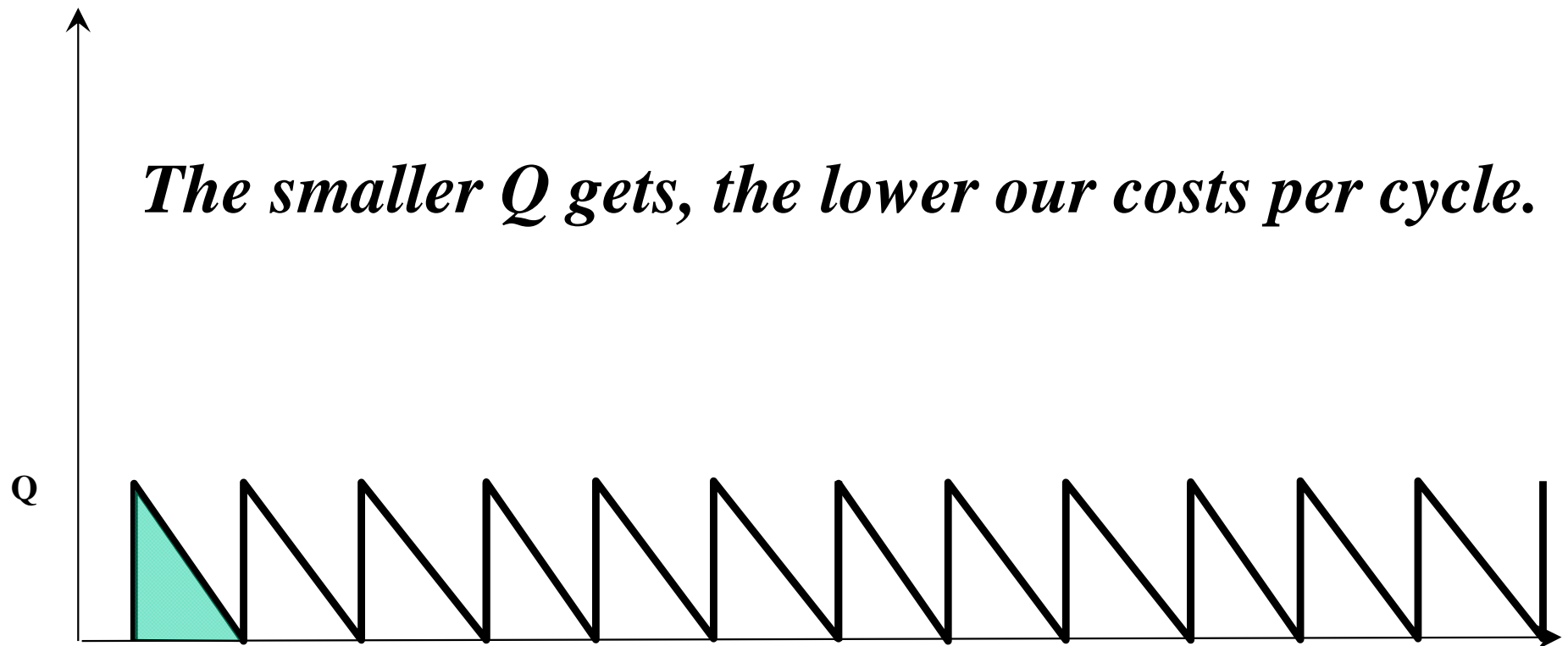
The economic order quantity



The economic order quantity



The economic order quantity



The economic order quantity

Need to minimize cost per unit time, not cost per order interval.

So, we want to solve:

$$\begin{aligned}\min C^\tau(Q) &= \frac{C(Q)}{\tau} = \frac{C(Q)}{Q/\lambda} = \frac{\lambda C(Q)}{Q} \\ &= \frac{\lambda}{Q} \left(K + c^p Q + c^h \left(\frac{Q^2}{2\lambda} \right) \right) \\ &= \frac{\lambda K}{Q} + \lambda c^p + c^h \frac{Q}{2}\end{aligned}$$

Differentiating with respect to Q and setting to 0:

$$\frac{dC^\tau(Q)}{dQ} = -\frac{\lambda K}{Q^2} + \frac{c^h}{2} = 0$$

The economic order quantity

Finally, solving for Q gives us:

$$Q = \sqrt{\frac{2K\lambda}{c^h}} = \text{The Economic Order Quantity (EOQ)}$$

Also called the Economic Lot Size.

Properties of the optimal solution:

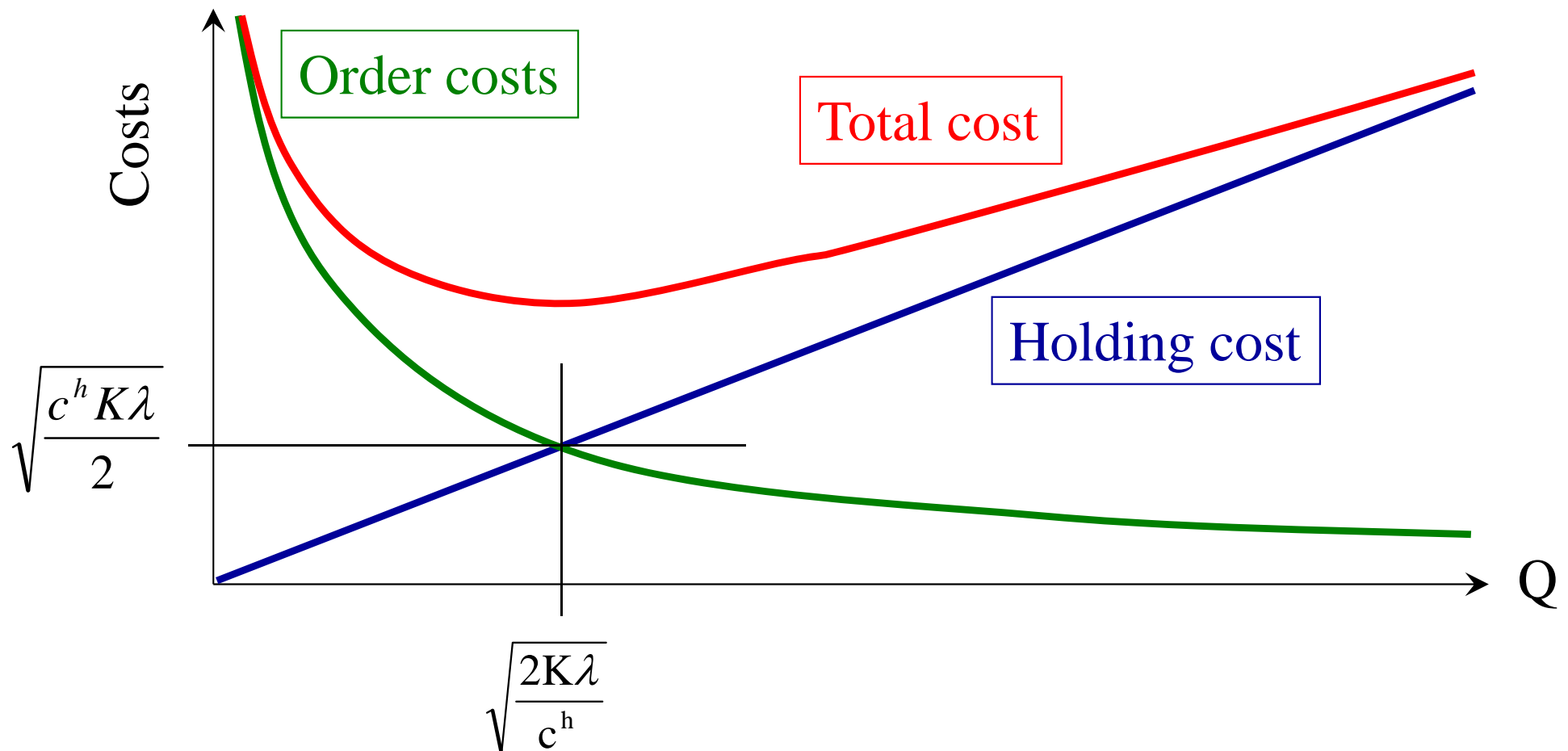
1. Purchase costs do not enter the equation (why?)
2. Order costs per unit time = holding costs per unit time:

$$\text{Order costs per unit time} = \frac{\lambda K}{Q} = \frac{\lambda K}{\sqrt{\frac{2K\lambda}{c^h}}} = \sqrt{\frac{c^h K \lambda}{2}}$$

$$\text{Holding costs per unit time} = c^h \frac{Q}{2} = \frac{c^h}{2} \sqrt{\frac{2K\lambda}{c^h}} = \sqrt{\frac{c^h K \lambda}{2}}$$

The economic order quantity

- The average cost function:



The economic order quantity

■ Sensitivity analysis



Outline

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An inventory game

■ The basics:

- » Random demand (uniform between 0 and 10)
- » Demands are not revealed until an order is entered.
- » Cost parameters:

Order cost	25
Purchase	10
Holding	1
Stockout	5

An inventory game

Player: Joe

Time	Inventc	Order	Demar	Ending	Sold	Lost	Order	Purcha	Holdin	Stocko	Total	Cum	
0	20	0	2	18	2	0	0	0	18	0	18	18	
1	18	0	1	17	1	0	0	0	17	0	17	35	
2	17	0	8	9	8	0	0	0	9	0	9	44	
3	9	0	3	6	3	0	0	0	6	0	6	50	
4	6	0	2	4	2	0	0	0	4	0	4	54	
5	4	4	6	2	6	0	25	40	2	0	67	121	
6	2	6	4	4	4	0	25	60	4	0	89	210	
7	4	4	5	3	5	0	25	40	3	0	68	278	
8	3	5	8	0	8	0	25	50	0	0	75	353	
9	0	8	6	2	6	0	25	80	2	0	107	460	
10	2	6	2	6	2	0	25	60	6	0	91	551	
11	6	2	3	5	3	0	25	20	5	0	50	601	
12	5	3	5	3	5	0	25	30	3	0	58	659	
13	3	5	5	3	5	0	25	50	3	0	78	737	
14	3	5	0	8	0	0	25	50	8	0	83	820	
15	8	0	2	6	2	0	0	0	6	0	6	826	
16	6	2	9	0	8	1	25	20	0	5	50	876	
17	0	8	1	7	1	0	25	80	7	0	112	988	
18	7	1	6	2	6	0	25	10	2	0	37	1025	
19	2	5	4	3	4	0	25	50	3	0	78	1103	
20	3	4	8	0	7	1	25	40	0	5	70	1173	
							375	680	108	10	1173		40

An inventory game Player Jimmie:

Time	Inventc	Order	Demar	Ending	Sold	Lost	Order	Purcha	Holdin	Stocko	Total	Cum	
0	20	0	2	18	2	0	0	0	18	0	18	18	
1	18	0	1	17	1	0	0	0	17	0	17	35	
2	17	0	8	9	8	0	0	0	9	0	9	44	
3	9	0	3	6	3	0	0	0	6	0	6	50	
4	6	0	2	4	2	0	0	0	4	0	4	54	
5	4	20	6	18	6	0	25	200	18	0	243	297	
6	18	0	4	14	4	0	0	0	14	0	14	311	
7	14	0	5	9	5	0	0	0	9	0	9	320	
8	9	0	8	1	8	0	0	0	1	0	1	321	
9	1	20	6	15	6	0	25	200	15	0	240	561	
10	15	0	2	13	2	0	0	0	13	0	13	574	
11	13	0	3	10	3	0	0	0	10	0	10	584	
12	10	0	5	5	5	0	0	0	5	0	5	589	
13	5	20	5	20	5	0	25	200	20	0	245	834	
14	20	0	0	20	0	0	0	0	20	0	20	854	
15	20	0	2	18	2	0	0	0	18	0	18	872	
16	18	0	9	9	9	0	0	0	9	0	9	881	
17	9	0	1	8	1	0	0	0	8	0	8	889	
18	8	0	6	2	6	0	0	0	2	0	2	891	
19	2	10	4	8	4	0	25	100	8	0	133	1024	
20	8	0	8	0	8	0	0	0	0	0	0	1024	
							100	700	224	0	1024		41

An inventory game Player:

Time	Inventc	Order	Demar	Ending	Sold	Lost	Order	Purcha	Holdin	Stocko	Total	Cum	
0	20	-1	0	20	-1	0	0	0	20	0	20		
1	20	-1	0	20	-1	0	0	0	20	0	20		
2	20	-1	0	20	-1	0	0	0	20	0	20		
3	20	-1	0	20	-1	0	0	0	20	0	20		
4	20	-1	0	20	-1	0	0	0	20	0	20		
5	20	-1	0	20	-1	0	0	0	20	0	20		
6	20	-1	0	20	-1	0	0	0	20	0	20		
7	20	-1	0	20	-1	0	0	0	20	0	20		
8	20	-1	0	20	-1	0	0	0	20	0	20		
9	20	-1	0	20	-1	0	0	0	20	0	20		
10	20	-1	0	20	-1	0	0	0	20	0	20		
11	20	-1	0	20	-1	0	0	0	20	0	20		
12	20	-1	0	20	-1	0	0	0	20	0	20		
13	20	-1	0	20	-1	0	0	0	20	0	20		
14	20	-1	0	20	-1	0	0	0	20	0	20		
15	20	-1	0	20	-1	0	0	0	20	0	20		
16	20	-1	0	20	-1	0	0	0	20	0	20		
17	20	-1	0	20	-1	0	0	0	20	0	20		
18	20	-1	0	20	-1	0	0	0	20	0	20		
19	20	-1	0	20	-1	0	0	0	20	0	20		
20	20	-1	0	20	-1	0	0	0	20	0	20		
							© 2013 W.B. Powell	0	420	0	420		42

An inventory game Player:

Time	Inventc	Order	Demar	Ending	Sold	Lost	Order	Purcha	Holdin	Stocko	Total	Cum	
0	20	-1	0	20	-1	0	0	0	20	0	20		
1	20	-1	0	20	-1	0	0	0	20	0	20		
2	20	-1	0	20	-1	0	0	0	20	0	20		
3	20	-1	0	20	-1	0	0	0	20	0	20		
4	20	-1	0	20	-1	0	0	0	20	0	20		
5	20	-1	0	20	-1	0	0	0	20	0	20		
6	20	-1	0	20	-1	0	0	0	20	0	20		
7	20	-1	0	20	-1	0	0	0	20	0	20		
8	20	-1	0	20	-1	0	0	0	20	0	20		
9	20	-1	0	20	-1	0	0	0	20	0	20		
10	20	-1	0	20	-1	0	0	0	20	0	20		
11	20	-1	0	20	-1	0	0	0	20	0	20		
12	20	-1	0	20	-1	0	0	0	20	0	20		
13	20	-1	0	20	-1	0	0	0	20	0	20		
14	20	-1	0	20	-1	0	0	0	20	0	20		
15	20	-1	0	20	-1	0	0	0	20	0	20		
16	20	-1	0	20	-1	0	0	0	20	0	20		
17	20	-1	0	20	-1	0	0	0	20	0	20		
18	20	-1	0	20	-1	0	0	0	20	0	20		
19	20	-1	0	20	-1	0	0	0	20	0	20		
20	20	-1	0	20	-1	0	0	0	20	0	20		
							© 2013 W.B. Powell	0	420	0	420		43

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Math programming formulation

■ What do we do when the demands are nonstationary?

» D_t = the *forecasted* demand for time period t , $0 \leq t < T$.
We are going to use a *point* estimate of the demand, which produces a *deterministic* model.

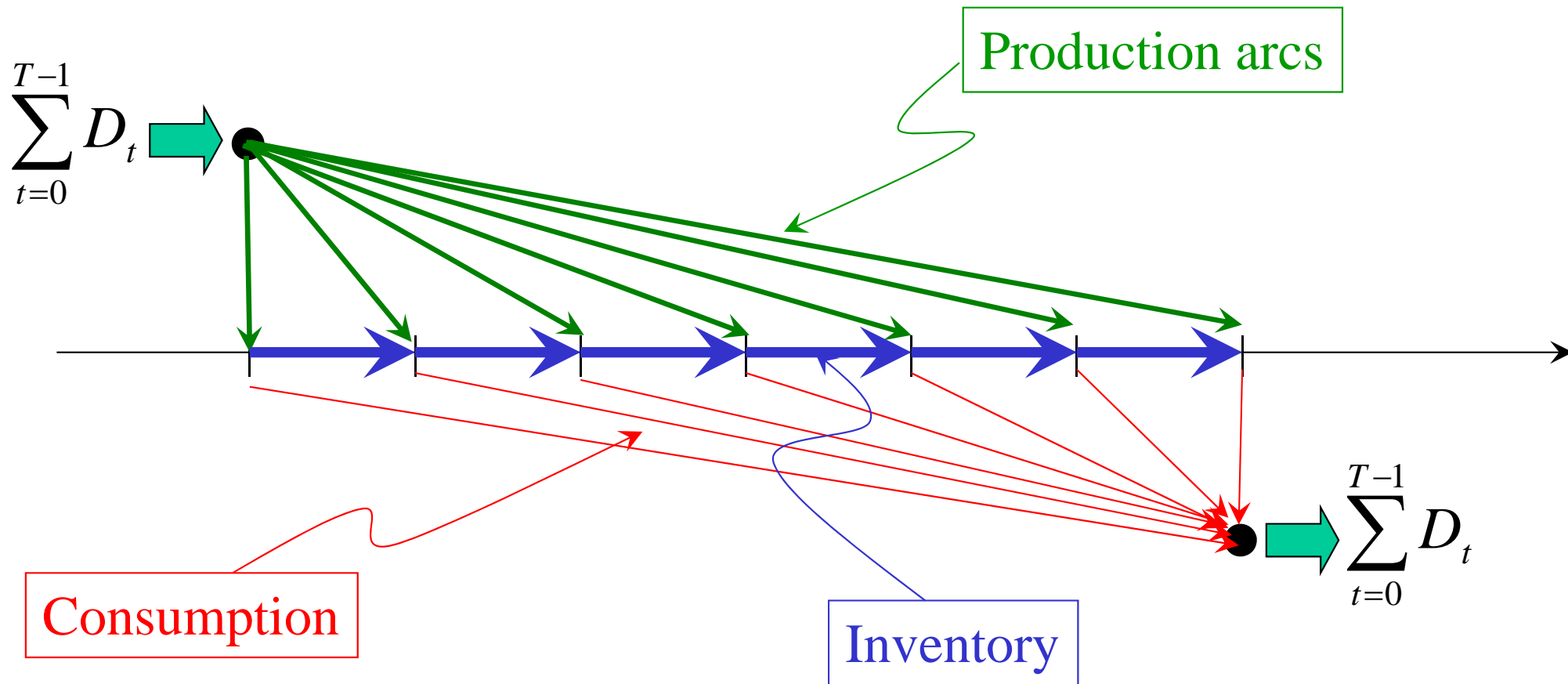
» T = the *planning horizon*.

■ Objective function:

$$\min_{x_t, y_t} \sum_{t=0}^{T-1} c_t (R_t, x_t, y_t)$$

Math programming formulation

- The optimization problem can be visualized as a network:



Math programming formulation

■ Integer programming formulation:

Notation:

Activity variables:

R_t = inventory at beginning of period t

D_t = demand during period starting at t

Parameters:

c^h = Unit holding cost per time period

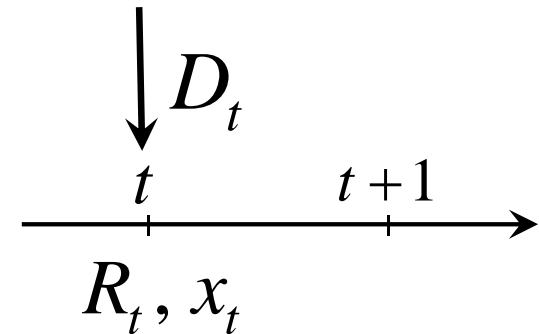
c^p = Unit purchase cost

K = Fixed order cost

Decision variables:

x_t = amount ordered in period t

$$y_t = \begin{cases} 1 & x_t > 0 \\ 0 & x_t = 0 \end{cases}$$



Math programming formulation

Objective function:

$$\min_{x,y} \sum_{t=0}^{T-1} Ky_t + c_t^p x_t + c^h (R_t + x_t - D_t)$$

subject to:

$$R_{t+1} = R_t + x_t - D_t$$

$$x_t \leq My_t \quad (\text{M} = \text{big number})$$

$$x_t \geq D_t - R_t$$

$$x_t \geq 0$$

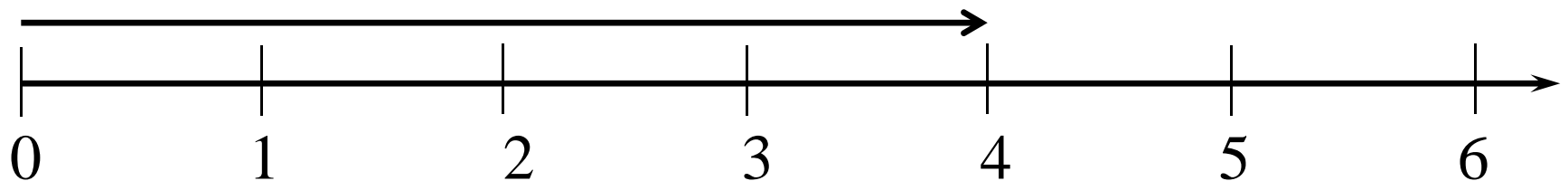
$$y_t = (0,1)$$

This is an integer programming problem, which can be solved using commercial solvers such as Cplex and Gurobi.

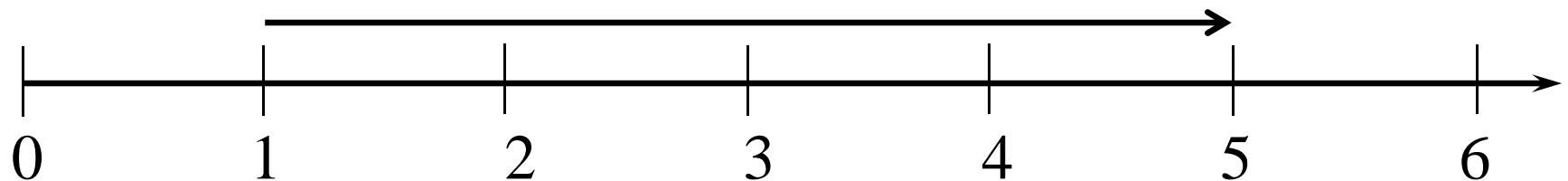
Math programming formulation

- We implement our math programming formulation as a *rolling horizon procedure*

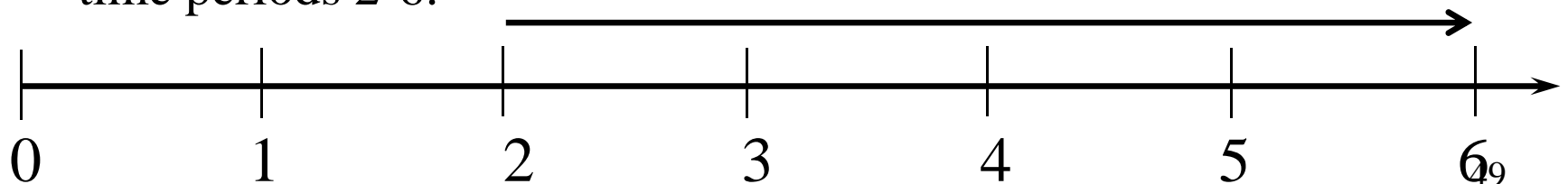
- » Optimize over 0-4, implement time 0



- » Roll to time 1, see new information, solve updated problem for time periods 1-5:



- » Roll to time 2, see new information, solve updated problem for time periods 2-6:



Math programming formulation

■ Rolling horizon procedures

» These are *deterministic approximations* of the problem over a *planning horizon* H

Objective function:

$$\min_{x_t, y_t} \sum_{t'=t}^{t+H} K y_{tt'} + c_t^p x_{tt'} + c^h (R_{tt'} + x_{tt'} - D_{tt'})$$

where $x_t = (x_{tt'})_{t'=t, \dots, t+H}$, $y_t = (y_{tt'})_{t'=t, \dots, t+H}$

subject to:

$$R_{t, t'+1} = R_{tt'} + x_{tt'} - D_{tt'}$$

$$x_{tt'} \leq M y_{tt'} \quad (M = \text{big number})$$

$$x_{tt'} \geq D_{tt'} - R_{tt'}$$

$$x_{tt'} \geq 0$$

$$y_{tt'} = (0, 1)$$

Heuristics

■ Silver-Meal heuristic (Least average cost)

Let $C(s)$ = average cost per unit time if we order over the next s time periods

$$= \frac{1}{s} \left(K + c^h \sum_{t=0}^{s-1} tD_t \right)$$

Calculate $C(1), C(2), \dots, C(s)$. Stop when $C(s+1) > C(s)$.

Set $T = s$. Order enough for the next T time periods.

» *One of the best known and most widely used heuristics in supply chain management.*

Heuristics

■ Least unit cost

Let $C(s)$ = average cost per unit produced if we order over the next s time periods

$$= \frac{\left(K + c^h \sum_{t=0}^{s-1} tD_t \right)}{\left(\sum_{t=0}^{s-1} D_t \right)}$$

Calculate $C(1), C(2), \dots, C(s)$. Stop when $C(s+1) > C(s)$.

Set $T = s$.

Least unit cost reflects the way managers are actually measured.

No - one is measured in terms of \$/day (why not???)

Wagner-Whitin algorithm

■ Example: Seasonal TV demand

» Parameters:

Setup cost	8
Order cost	1
Holding cost	1.3

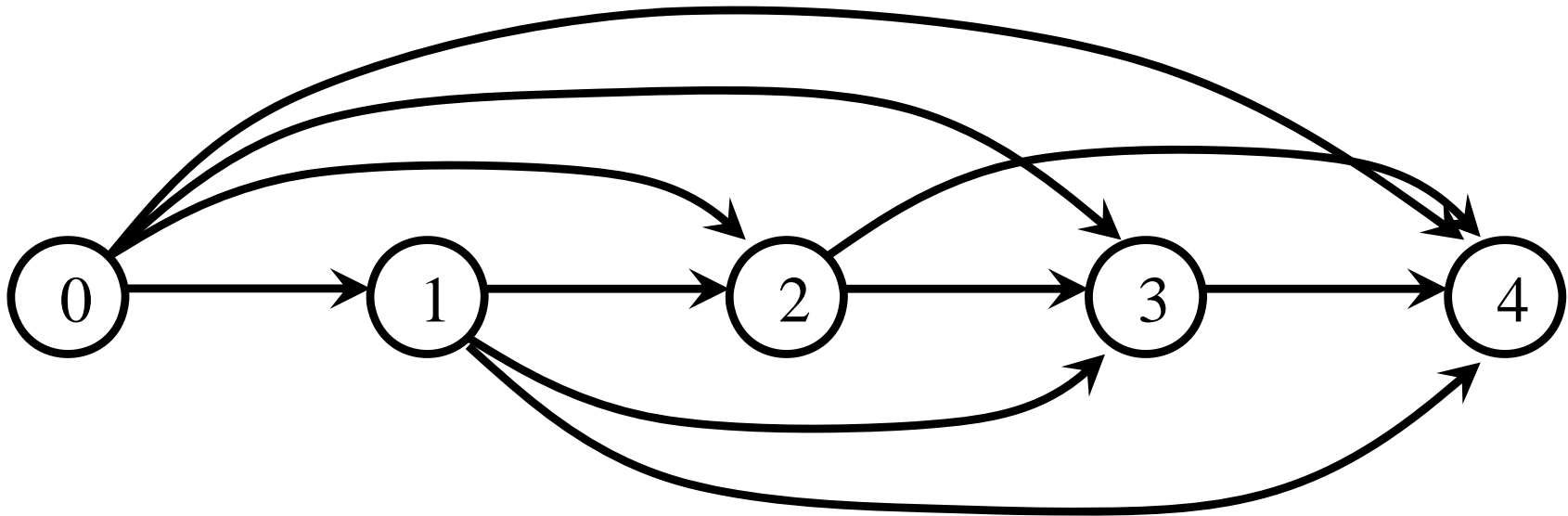
» Demands (in 1000's):

Time	Demand
0	1
1	3
2	5
3	2

Wagner-Whitin algorithm

■ A network representation:

- » The set of decisions represents a shortest path problem over a specialized network:

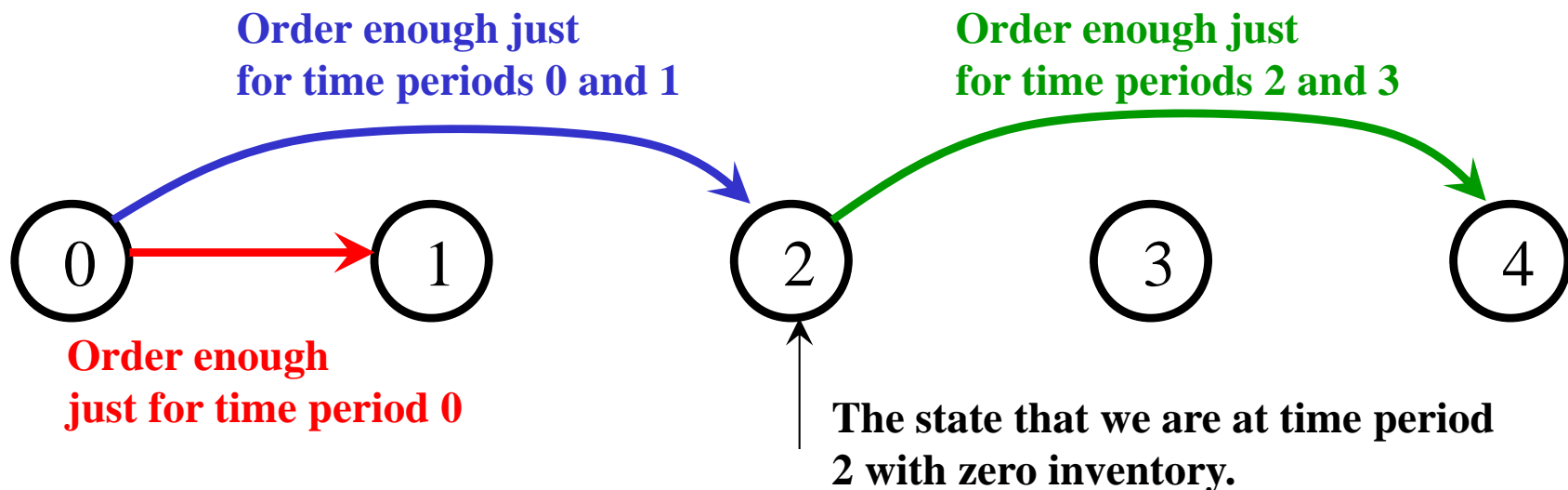


- » The cost on each arc is the cost of the decision, including order costs and all holding costs.
- » What does the optimal solution look like?

Wagner-Whitin algorithm

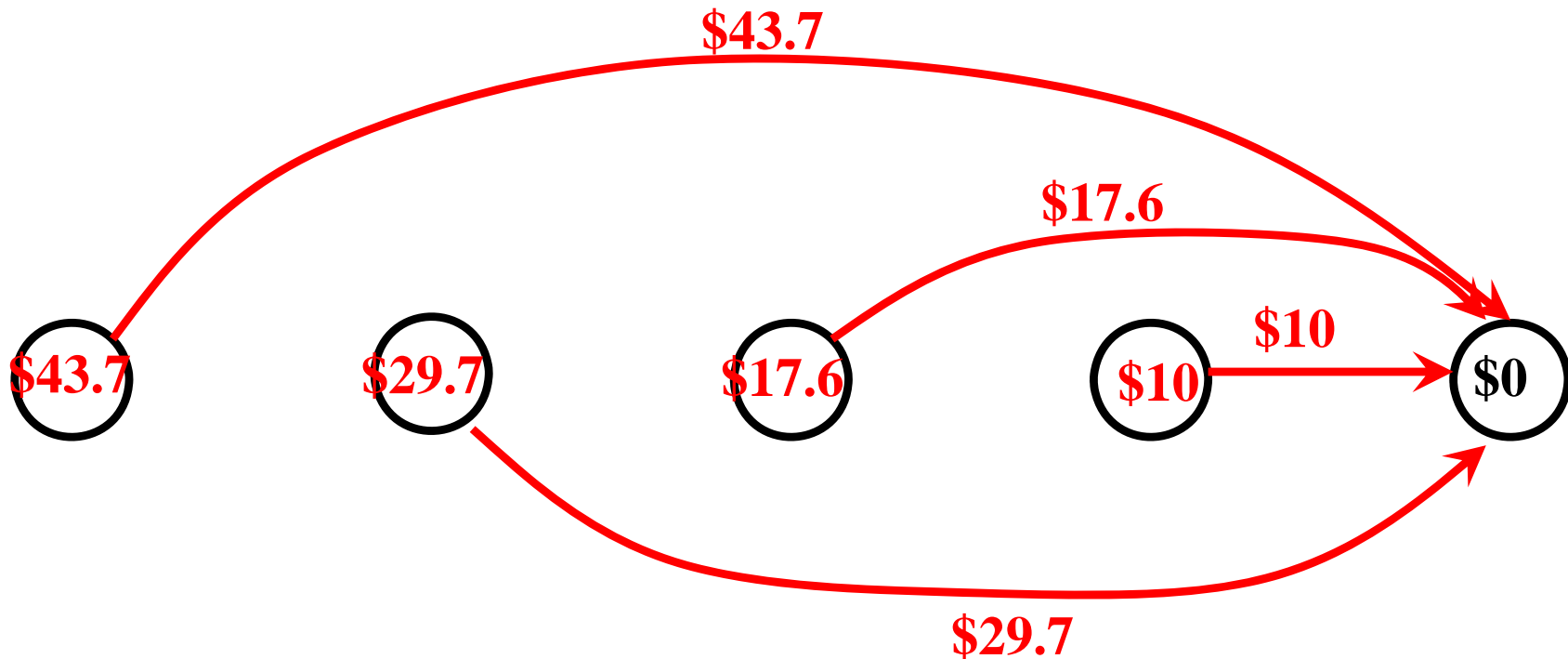
■ Properties of optimal solution

- » We only need to make decisions when the inventory is zero.
- » This means our decision variable is not the quantity, but the number of time periods into the future that we need to cover.



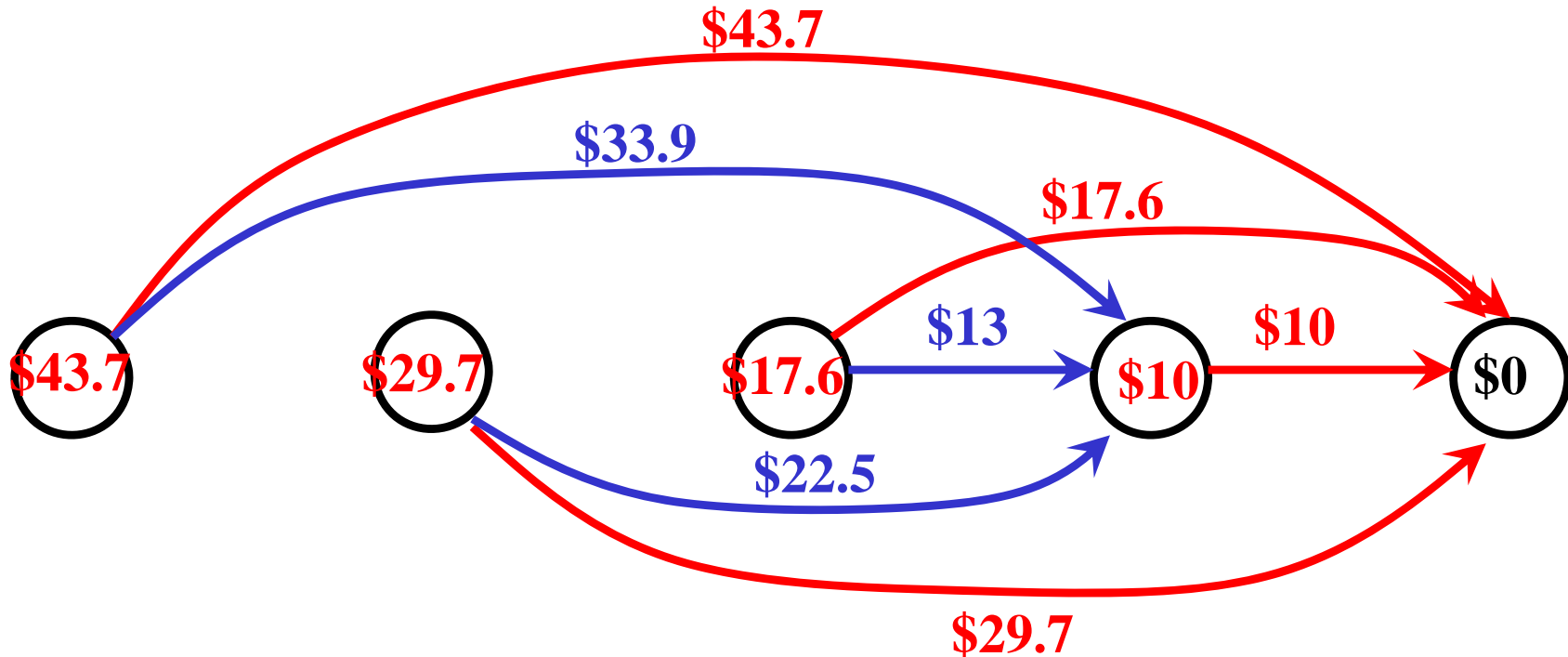
Wagner-Whitin algorithm

■ Start with the final node:



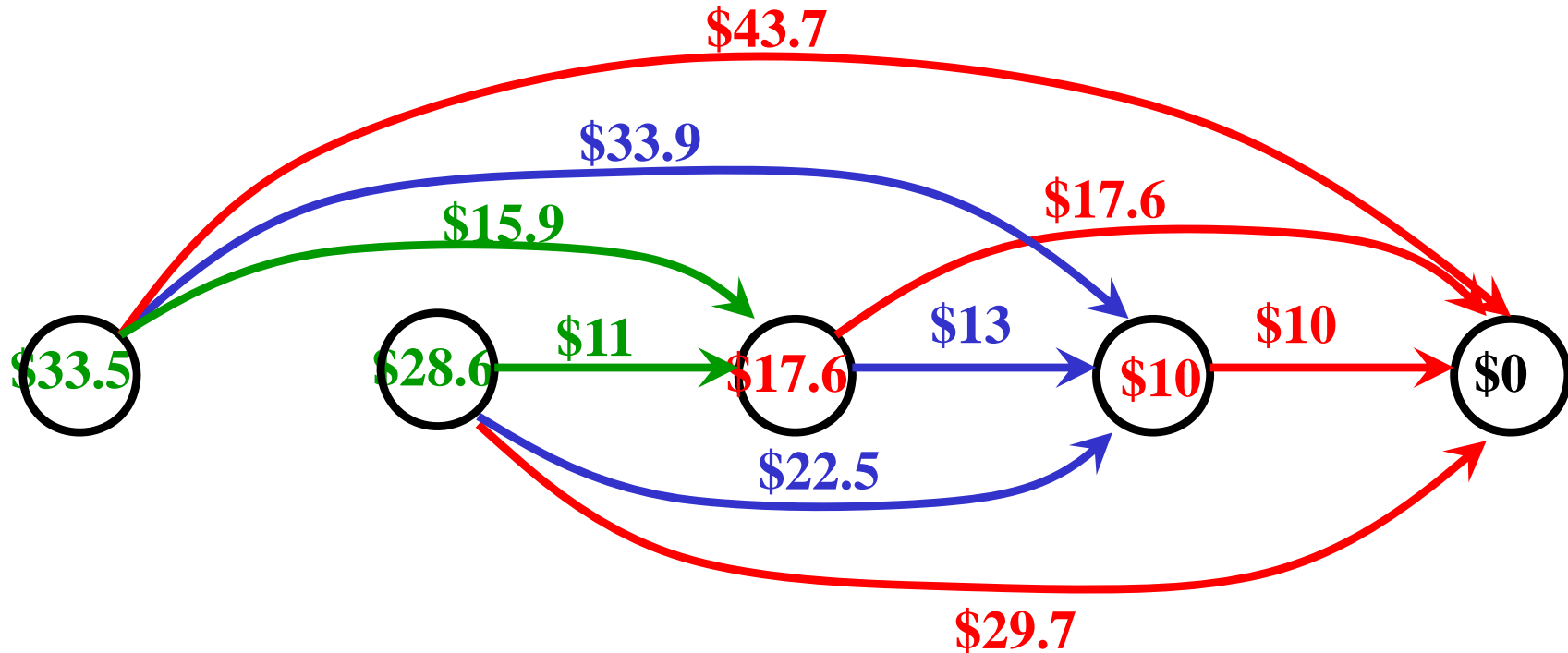
Wagner-Whitin algorithm

■ Links into time period 3:



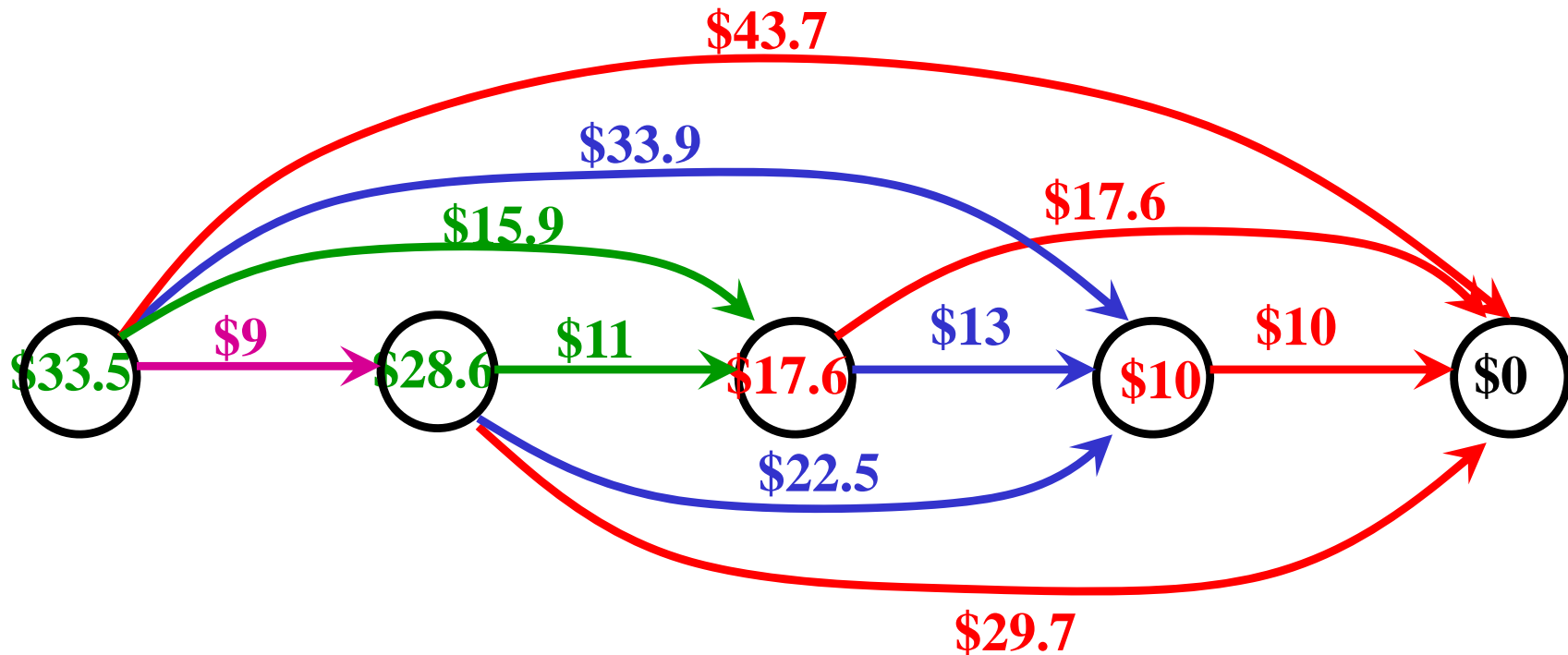
Wagner-Whitin algorithm

■ Links into time period 2:



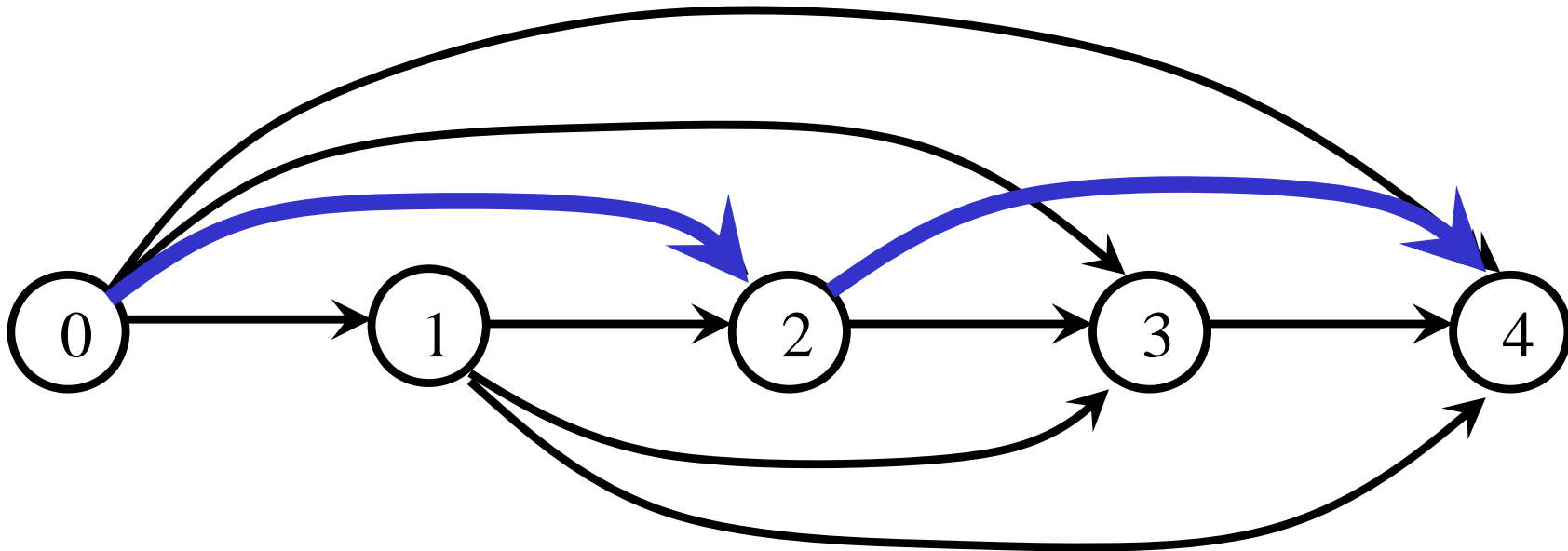
Wagner-Whitin algorithm

■ Links into time period 1:



Wagner-Whitin algorithm

- Finally, we walk forward in time:



We use the values computed in the backward pass to walk forward and compute decisions.

Wagner-Whitin algorithm

■ Strengths:

- » Very fast
- » Handles very general cost functions
 - You can use virtually any shape order cost function.
- » Handles time-dependent data (e.g. seasonal data, day of week effects or hour of day patterns).
- » Handles forecasts of the future (which is a form of time-dependency).

Wagner-Whitin algorithm

■ Limitations of this model:

» Assumes demands are deterministic!!!

- “Optimal” solution is not really optimal.
- Drives inventories to zero, which will create stockouts.
- Have to reoptimize as forecasted demands change.

» Limitations:

- Computationally demanding when you have to solve 100,000 problems (Wal-Mart!).
- Solutions are not “obviously” better than good heuristics under realistic conditions.
- Gets complicated if you have multiple items and joint capacity constraints (need to use integer programming formulation)

» But:

- Serves as a useful subproblem in the context of larger applications.
- Highlights behavior of the problem.

Wagner-Whitin algorithm

■ Important generalizations:

- » Upper bounds on order quantities
 - What if we cannot order more than u_t in time period t ?
- » Upper bounds on production and multiple items:
 - This is the problem that actually arises in practice.
 - Called the “capacitated multi-item lot sizing problem.”
- » Limit on total production time, in the presence of setup times.
 - The literature on setup times is very sparse.